



Introduction to Computational Logic, SS 2006: Solution for assignment 11

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Exercise 11.1 (Inconsistence) a) Claim: $QL \not\vdash 0 = 1$

Proof: By Contradiction. Let $QL \vdash 0 = 1$. Then $\mathcal{A}0 = \mathcal{A}1$ in every model of QL (by soundness). Contradiction.

b)

$$\begin{array}{lll} \text{"}\neg\text{"}: & 0 = \bar{0} & x = \bar{x} \\ & = 1 & \text{Negation} \end{array}$$

$$\begin{array}{lll} \text{"}\vdash\text{"}: & x = x \wedge 1 & \text{Identity} \\ & = x \wedge 0 & 0 = 1 \\ & = 0 & \text{Dominance} \\ & = \bar{x} \wedge 0 & \text{Dominance} \\ & = \bar{x} \wedge 1 & 0 = 1 \\ & = \bar{x} & \text{Identity} \end{array}$$

c)

$$\text{"}\neg\text{"}: \quad 0 = 1 \quad x = y$$

$$\begin{array}{lll} \text{"}\vdash\text{"}: & x = x \wedge 1 & \text{Identity} \\ & = x \wedge 0 & 0 = 1 \\ & = 0 & \text{Dominance} \\ & = y \wedge 0 & \text{Dominance} \\ & = y \wedge 1 & 0 = 1 \\ & = y & \text{Identity} \end{array}$$

Exercise 11.2 (Duality) $e = (a = 0)$ where a is a parameter.

Exercise 11.3 (\exists Defined)

a)

$$\begin{array}{ll}
 \exists x.0 = \overline{\forall x.(\lambda x.0)x} & \exists f = \overline{\forall x.\overline{fx}} \\
 = \overline{\forall x.0} & \beta \\
 = \overline{\forall x.1} & \text{Taut} \\
 = \bar{1} & \forall 1 \\
 = 0 & \text{Taut}
 \end{array}$$

b)

$$\begin{array}{ll}
 fx \rightarrow \exists f = fx \rightarrow \overline{\forall x.\overline{fx}} & \exists f = \overline{\forall x.\overline{fx}} \\
 = \forall x.\overline{fx} \rightarrow \overline{fx} & \text{Taut} \\
 = \forall x.\overline{fx} \rightarrow (\lambda x.\overline{fx})x & \beta \\
 = 1 & \forall I
 \end{array}$$

c) " \rightarrow ": follows from a) and b).

" \vdash ": We have to show $QL \vdash \exists f = \overline{\forall x.\overline{fx}}$.

$$\begin{array}{ll}
 \overline{\forall x.\overline{fx}} = \exists x.\overline{\overline{fx}} & \text{dM} \\
 = \exists x.\overline{fx} & \text{Taut} \\
 = \exists f & \eta
 \end{array}$$

Exercise 11.4 (Pull Laws \rightarrow)

a)

$$\begin{array}{ll}
 q \rightarrow \exists f = \bar{q} \vee \exists f & \text{Taut} \\
 = \exists x.\bar{q} \vee fx & \exists \vee \\
 = \exists x.q \rightarrow fx & \text{Taut}
 \end{array}$$

b)

$$\begin{array}{ll}
 q \rightarrow \forall f = \bar{q} \vee \forall f & \text{Taut} \\
 = \forall x.\bar{q} \vee fx & \forall \vee \\
 = \forall x.q \rightarrow fx & \text{Taut}
 \end{array}$$

c)

$$\begin{aligned}\exists f \rightarrow q &= \overline{\exists f} \vee q && \text{Taut} \\ &= (\forall x.\overline{fx}) \vee q && \text{dM} \\ &= \forall x.\overline{fx} \vee q && \forall\vee \\ &= \forall x.fx \rightarrow q && \text{Taut}\end{aligned}$$

d)

$$\begin{aligned}\forall f \rightarrow q &= \overline{\forall f} \vee q && \text{Taut} \\ &= (\exists x.\overline{fx}) \vee q && \text{dM} \\ &= \exists x.\overline{fx} \vee q && \exists\vee \\ &= \exists x.fx \rightarrow q && \text{Taut}\end{aligned}$$

e)

$$\begin{aligned}\forall f \rightarrow \exists g &= \overline{\forall f} \vee \exists g && \text{Taut} \\ &= (\exists x.\overline{fx}) \vee \exists g && \eta, \text{dM} \\ &= \exists x.\overline{fx} \vee gx && \exists\vee', \beta \\ &= \exists x.fx \rightarrow gx && \text{Taut}\end{aligned}$$

Exercise 11.5 (Backward Proofs)

a)

$$\begin{aligned}\forall f \rightarrow \exists f &= \exists x.fx \rightarrow fx && \text{Pull } \rightarrow \\ &= \exists x.1 && \text{Taut} \\ &= 1 && \exists E\end{aligned}$$

b)

$$\begin{aligned}\exists x.fx \rightarrow \forall f &= \forall f \rightarrow \forall f && \text{Pull } \rightarrow \\ &= 1 && \text{Taut}\end{aligned}$$

c)

$$\begin{aligned} & \forall f \vee \forall g \rightarrow \forall x. fx \vee gx \\ = & \forall x. \forall f \vee \forall g \rightarrow fx \vee gx && \text{Pull } \rightarrow \\ \vdash & \forall f \vee \forall g \rightarrow fx \vee gx && \text{Gen } \forall \\ = & (\forall yz. fy \vee gz) \rightarrow fx \vee gx && \forall \vee (2x) \\ = & \exists yz. fy \vee gz \rightarrow fx \vee gx && \text{Pull } \rightarrow (2x) \\ \neg & fx \vee gx \rightarrow fx \vee gx && \text{Gen } \exists, y = x, z = x \\ = & 1 && \text{Taut} \end{aligned}$$

d)

$$\begin{aligned} & \forall f \rightarrow \forall x. fx \vee gx \\ = & \forall x. \forall f \rightarrow fx \vee gx && \text{Pull } \rightarrow \\ = & \forall x. (\forall f \rightarrow fx) \vee gx && \text{Taut} \\ = & \forall x. 1 \vee gx && \forall I \\ = & \forall x. 1 && \text{Taut} \\ = & 1 && \forall 1 \end{aligned}$$

e)

$$\begin{aligned} & (\exists f \rightarrow \forall h) \rightarrow \forall x. fx \rightarrow hx \\ \vdash & (\exists f \rightarrow \forall h) \rightarrow fx \rightarrow hx && \text{Pull } \rightarrow, \text{Gen } \forall \\ = & (\forall y \forall z. fy \rightarrow hz) \rightarrow fx \rightarrow hx && \text{Pull } \rightarrow \\ = & \exists y \exists z. (fy \rightarrow hz) \rightarrow fx \rightarrow hx && \text{Pull } \rightarrow \\ \neg & (fx \rightarrow hx) \rightarrow fx \rightarrow hx && \text{Gen } \exists, y = x, z = x \\ = & 1 && \text{Taut} \end{aligned}$$

f)

$$\begin{aligned} & (\forall x. fx \vee hx) \rightarrow \forall f \vee \exists h \\ = & \forall y \exists z. (\forall x. fx \vee hx) \rightarrow fy \vee hz && \forall \vee, \exists \vee, \text{Pull} \rightarrow \\ \neg & (\forall x. fx \vee hx) \rightarrow fy \vee hy && \text{Gen}\forall, \text{Gen}\exists, z = y \\ = & \exists x. (fx \vee hx) \rightarrow fy \vee hy && \text{Pull} \rightarrow \\ \neg & (fy \vee hy) \rightarrow fy \vee hy && \text{Gen}\exists, x = y \\ = & 1 && \text{Taut} \end{aligned}$$

g)

$$\begin{aligned} & (\exists x \forall y. fxy) \rightarrow \forall y \exists x. fxy \\ & (\exists x \forall y. fxy) \rightarrow \exists x. fxy && \text{Pull} \rightarrow, \text{Gen}\forall \\ \forall x \exists y'. fxy' \rightarrow \exists x. fxy && \text{Pull} \rightarrow \\ fxy \rightarrow \exists x. fxy && \text{Gen}\forall, \text{Gen}\exists, y' = y \\ \exists x'. fxy \rightarrow fx'y && \text{Pull} \rightarrow \\ fxy \rightarrow fxy && \text{Gen}\exists, x' = x \\ & 1 && \text{Taut} \end{aligned}$$

h)

$$\begin{aligned} & \forall x. fx \rightarrow \exists f \\ \forall x \exists x'. fx \rightarrow fx' && \text{Pull} \rightarrow \\ fx \rightarrow fx && \text{Gen}\forall, \text{Gen}\exists, x' = x \\ & 1 && \text{Taut} \end{aligned}$$

i)

$$\begin{aligned} & (\forall x. fx \rightarrow gx) \rightarrow \forall f \rightarrow \forall g \\ & (\forall x. fx \rightarrow gx) \rightarrow \forall f \rightarrow gx && \text{Pull} \rightarrow, \text{Gen}\forall \\ \exists x'. (\forall x. fx \rightarrow gx) \rightarrow fx' \rightarrow gx && \text{Pull} \rightarrow \\ & (\forall x. fx \rightarrow gx) \rightarrow fx \rightarrow gx && \text{Gen}\exists, x' = x \\ \exists x'. (fx' \rightarrow gx') \rightarrow fx \rightarrow gx && \text{Pull} \rightarrow \\ & (fx \rightarrow gx) \rightarrow fx \rightarrow gx && \text{Gen}\exists, x' = x \\ & 1 && \text{Taut} \end{aligned}$$

j)

$(\forall x. fx \rightarrow gx) \rightarrow \exists f \rightarrow \exists g$	
$(\forall x. fx \rightarrow gx) \rightarrow fx \rightarrow \exists g$	Pull \rightarrow , Gen \forall
$\exists x'. (\forall x. fx \rightarrow gx) \rightarrow fx \rightarrow gx'$	Pull \rightarrow
$(\forall x. fx \rightarrow gx) \rightarrow fx \rightarrow gx$	Gen \exists , $x' = x$
remaining steps as in 11.5 i)	

k)

$(\forall x. fx \rightarrow \overline{gx}) \rightarrow \forall y. gy \rightarrow \overline{fy}$	
$(\forall x. fx \rightarrow \overline{gx}) \rightarrow gy \rightarrow \overline{fy}$	Pull \rightarrow , Gen \forall
$\exists x. (fx \rightarrow \overline{gx}) \rightarrow gy \rightarrow \overline{fy}$	Pull \rightarrow
$(fy \rightarrow \overline{gy}) \rightarrow gy \rightarrow \overline{fy}$	Gen \exists , $x = y$
1	Taut

Exercise 11.6 (Counterexample) Let $QL \vdash \forall f \vee \forall g = \forall x. fx \vee gx$. Then

$QL \vdash (\forall x.x) \vee (\forall x.\bar{x}) = \forall x. (\lambda x.x)x \vee (\lambda x.\bar{x})x$	Generativity
$QL \vdash (\forall x.x) \wedge 0 \vee (\forall x.\bar{x}) \wedge \bar{1} = \forall x. x \vee \bar{x}$	$\forall I$, GR, β
$QL \vdash 0 = \forall x.1$	Taut
$QL \vdash 0 = 1$	$\forall 1$

Contradiction!

Exercise 11.7 (Deductivity)

$s = a \wedge f1$
 $t = fa$

Property 2) is satisfied since

$$a \wedge f1 = 1 \stackrel{BA}{\vdash} \{a = 1, f1 = 1\} \stackrel{BA}{\vdash} fa = 1 \quad \text{by And}$$

Property 3) is satisfied due to the following reason.

First, $BA \vdash a \wedge f1 \rightarrow fa = \bar{a} \vee \overline{f1} \vee fa$.

Now consider a powerset algebra \mathcal{A} over the set $\{1, 2\}$ s.t.

$$\mathcal{A}a = \{1\}$$

$$\mathcal{A}f\{1, 2\} = \{1, 2\}$$

$$\mathcal{A}f\{1\} = \emptyset$$

BA $\not\models a \wedge f1 \rightarrow fa = 1$ now follows from the fact that \mathcal{A} does not satisfy $\bar{a} \vee \overline{f1} \vee fa = 1$.