



Assignment 9 Introduction to Computational Logic, SS 2006

Prof. Dr. Gert Smolka, Dipl.-Inform. Mathias Möhl
www.ps.uni-sb.de/courses/cl-ss06/

Exercise 9.1 Find an entailment relation on \mathbb{N} that is not compact.

Exercise 9.2 Let \vdash be an entailment relation on X . Prove:

$$A \vdash A' \iff \forall x \in X: A' \vdash x \implies A \vdash x$$

Exercise 9.3 Consider the inference system

$$\{(\{x, y\}, x + y) \mid x, y \in \mathbb{N}\}$$

Prove that for every $x \in S[\{3, 15, 21\}]$ there exists an $n \in \mathbb{N}$ such that $x = 3n$.

Exercise 9.4 (Incompleteness) Let Nat be a categorical specification of the Booleans and the natural numbers that axiomatizes the constants

$$\begin{aligned} 0, 1: B \\ \sigma: N \rightarrow N \\ o: N \end{aligned}$$

with their usual meaning. Find A and e such that $\text{Nat}, A \models e$ and $\text{Nat}, A \not\models e$.

Hints: Choose an infinite A and exploit the compactness of deductive entailment. Use a constant $f: N \rightarrow B$ that does not appear in Nat . Choose e such that it requires f to represent an infinite set.

Exercise 9.5 (Termination) Let PL be a specification that axiomatizes the constants

$$\begin{aligned} 0, 1: B \\ \wedge, \rightarrow: B \rightarrow B \rightarrow B \\ \forall, \exists: (V \rightarrow B) \rightarrow B \end{aligned}$$

with their usual meaning once V is fixed.

- a) Let $r: V \rightarrow V \rightarrow B$ be a constant that doesn't occur in PL. Find an equation ter such that a model of PL is a model of $\text{PL} \cup \{\text{ter}\}$ if and only if it interprets r with a terminating relation.
- b) Find equations A such that
- (i) $\text{PL}, \text{ter}, A \models 0 = 1$
 - (ii) $\text{PL}, \text{ter}, A \not\models 0 = 1$
- Hint: Use pairwise distinct constants a_1, a_2, \dots that don't occur in PL and exploit the compactness of deductive entailment.

Exercise 9.6 (Finiteness) Let PL be a specification that axiomatizes the constants

$$\begin{aligned} 0, 1: & B \\ \rightarrow: & B \rightarrow B \rightarrow B \\ \forall, \exists: & (V \rightarrow B) \rightarrow B \\ \doteq: & V \rightarrow V \rightarrow B \end{aligned}$$

with their usual meaning once V is fixed.

- a) Find an equation fin such that a model of PL is a model of $\text{PL} \cup \{\text{fin}\}$ if and only if it interprets V with a finite set.
- Hint: A set X is finite if and only if every injective function $X \rightarrow X$ is surjective.
- b) Find equations A such that
- (i) $\text{PL}, \text{fin}, A \models 0 = 1$
 - (ii) $\text{PL}, \text{fin}, A \not\models 0 = 1$

Exercise 9.7 (Transitive Closure) Let PL be a specification that axiomatizes the constants

$$\begin{aligned} 0, 1: & B \\ \wedge, \vee, \rightarrow: & B \rightarrow B \rightarrow B \\ \forall, \exists: & (V \rightarrow B) \rightarrow B \\ \forall: & ((V \rightarrow V \rightarrow B) \rightarrow B) \rightarrow B \end{aligned}$$

with their usual meaning once V is fixed. We use the notation $\text{Rel} := V \rightarrow V \rightarrow B$.

- a) Give a term trans of type $\text{Rel} \rightarrow B$ that describes a function that tests whether a relation is transitive.
- b) Give a term incl of type $\text{Rel} \rightarrow \text{Rel} \rightarrow B$ that describes a function that tests

$r \subseteq r'$.

- c) Give a term tc of type $Rel \rightarrow Rel \rightarrow B$ that describes a function that tests whether the second argument is the transitive closure of the first argument.
- d) Let $r, r' : Rel$ be constants that don't occur in PL. Show that

$$PL \cup \{r'xy = rxy \vee (\exists z. rxz \wedge r'zy)\}$$

has models where r' is not the transitive closure of r .

Exercise 9.8 Draw a locally confluent relation that is not confluent.

Hint: Two non-terminal and two terminal nodes suffice.

Exercise 9.9 Draw an infinite relation $\rightarrow \subseteq \mathbb{N}^2$ that is confluent, non-terminating, and where every $n \in \mathbb{N}$ has exactly one normal form.

Exercise 9.10 Draw a terminating relation \rightarrow such that $\exists x \forall n \in \mathbb{N} \exists y : x \rightarrow^n y$.

Exercise 9.11 Make sure that you understand the proofs of the following facts.

- a) \rightarrow semi-confluent $\Rightarrow \rightarrow$ Church-Rosser
- b) \rightarrow terminating and locally confluent $\Rightarrow \rightarrow$ confluent

Exercise 9.12 Let $\rightarrow \subseteq X^2$ be confluent and terminating. Prove the following statements:

- a) There is no $x \in X$ that has two normal forms.
- b) $\forall x, y \in X : x \leftrightarrow^* y \Rightarrow x, y$ have the same normal forms.

You may use the following facts:

- i) \rightarrow confluent $\Rightarrow \rightarrow$ Church-Rosser
- ii) \rightarrow terminating $\Rightarrow \forall x \in X : x$ has a normal form