

# Computer Architecture 1 - Übungsblatt 3

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# Aufgabe 1:

**Kosten:**

$$\begin{aligned}C_{hdec}(1) &= 0 \\C_{hdec}(n) &= C_{hdec}(n-1) + 2^{n-1} \cdot (C_{and} + C_{or}) \\&= C_{hdec}(n-1) + 2^{n-1} \cdot 2 \\&= C_{hdec}(n-1) + 2^n \\&= 2^n + 2^{n-1} + \dots + 2^2 + C_{hdec}(1) \\&= 2^n + 2^{n-1} + \dots + 2^2 + 0 \\&= 2^n + 2^{n-1} + \dots + 2^2 + 2^1 + 2^0 - 3 \\&= \left( \sum_{i=0}^n 2^i \right) - 3 \\&= (2^{n+1} - 1) - 3 \\&= 2^{n+1} - 4\end{aligned}$$

**Beweis:**

$$\begin{aligned}C_{hdec}(1) &= 0 = 2^{1+1} - 4 \\C_{hdec}(n) &= 2 \cdot 2^{n-1} + C_{hdec}(n-1) \\&\stackrel{IV}{=} 2^n + (2^n - 4) \\&= 2^{n+1} - 4\end{aligned}$$

□

**Tiefe:**

$$\begin{aligned}D_{hdec}(1) &= 0 \\D_{hdec}(n) &= D_{hdec}(n-1) + \max\{D_{and}, D_{or}\} \\&= D_{hdec}(n-1) + 1 \\&= D_{hdec}(n-2) + 1 + 1 \\&= \dots \\&= n - 1 + D_{hdec}(1) \\&= n - 1 + 0 \\&= n - 1\end{aligned}$$

Beweis:

$$\begin{aligned}D_{hdec}(1) &= 0 = 1 - 1 \\D_{hdec}(n) &= D_{hdec}(n - 1) + 1 \\&\stackrel{IV}{=} n - 2 + 1 \\&= n - 1\end{aligned}$$

□

## Aufgabe 2:

(a)

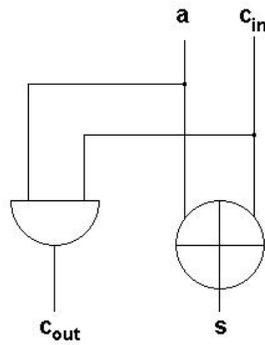


Abbildung 1: 1-bit Half-Adder

(b)

$$CCI(0) := (c_{in} = c_{out})$$

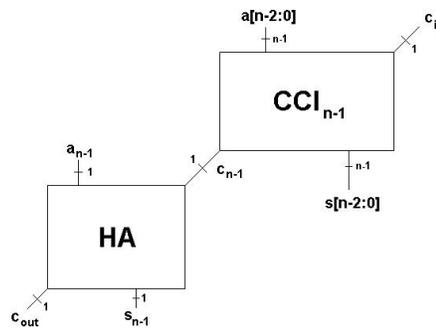


Abbildung 2: n-bit Half-Adder

(c)

Wertetabelle (\*):

$a$	$c_{in}$	$c_{out}$	$s$
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

Korrektheit des Half-Adder:

$$\langle a_0 \rangle + c_{in} \stackrel{(*)}{=} \langle c_{out}, s \rangle$$

Induktionsanfang:  $n = 0$

$$\langle a[-1 : 0] \rangle + c_{in} = 0 + c_{in} = c_{in} \stackrel{\text{n.Def.}}{=} c_{out} = \langle c_{out}, s[-1 : 0] \rangle$$

Induktionsschritt:  $n \leftrightarrow n + 1$

$$\begin{aligned} \langle a[n : 0] \rangle + c_{in} &= a_n \cdot 2^n + \langle a[n-1 : 0] \rangle + c_{in} \\ &\stackrel{\text{IA}}{=} a_n \cdot 2^n + \langle c_{n-1}, s[n-1 : 0] \rangle \\ &= a_n \cdot 2^n + c_{n-1} \cdot 2^n + \langle s[n-1 : 0] \rangle \\ &= (a_n + c_{n-1}) \cdot 2^n + \langle s[n-1 : 0] \rangle \\ &\stackrel{\text{HA}}{=} \langle c_n, s_n \rangle \cdot 2^n + \langle s[n-1 : 0] \rangle \\ &= \langle c_n, s[n : 0] \rangle \\ &= \langle c_{out}, s[n : 0] \rangle \end{aligned}$$

□

## Aufgabe 4:

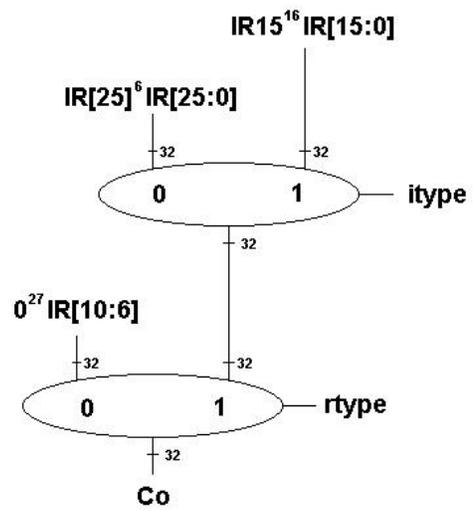


Abbildung 3: Co-Berechnung