

Problem 1 (Statecharts)

[20 points]

Draw a statechart modeling a traffic light controller that helps pedestrians cross a street.

Your model should contain two lights: A for the cars, and B for the pedestrians. Light A can be green, yellow, or red. Light B can be on, off, or blinking. Initially light A should be green, while light B should be off.

When light A is green and a **request** event is generated (meaning that a pedestrian wants to cross the street), then the following should happen:

Immediately: light A turns yellow;

After 10 seconds: light A turns red; light B turns on;

After 30 seconds: light B blinks;

After 1 minute: light A turns green; light B turns off.

Your model should include at least:

- states green, yellow, and red for light A ;
- states on, off, and blink for light B ;
- an external event **request**, generated when a pedestrian wants to cross the street;
- an internal event **tick**, periodically generated every second with the help of a timer.

Problem 2 (Petri nets)

[20 points]

Draw a colored Petri net implementing the function

$$\text{count}(a, X),$$

where

- a is an integer number;
- $X = \langle x_1, \dots, x_n \rangle$ is a nonempty multiset of integer numbers;
- $\text{count}(a, X)$ is the number of occurrences of a in X . For instance, $\text{count}(1, \langle 1, 1, 2 \rangle) = 2$ and $\text{count}(1, \langle 2, 3 \rangle) = 0$.

Your net should contain two “input” places A and B . A initially contains one token labeled by a . B initially contains n tokens labeled by the elements of X . All other places should initially contain no tokens. After a finite number of transitions, your net should enter a deadlock, and a given “output” place R should contain one token labeled by $\text{count}(a, X)$. If necessary, you can assign priorities to the transitions of your net. When more than one transition is activated, assume that only a transition with the highest priority can fire.

Problem 3 (Kahn process networks)

[20 points]

Draw a Kahn process network that generates the sequence

$$f(n) = n + 1, \quad \text{for } n \in \mathbb{N}.$$

You can use only the following processes:

- ① initially generates the number 1, then simply forwards its input;
- ⓓ duplicates a number;
- ⊕ sums two numbers;
- Ⓢ receives a number and outputs nothing; this is the sink for the sequence $f(n)$.

Problem 4 (VHDL)

[20 points]

Starting with the following VHDL declaration,

```
entity clb is
  port (a : in Bit;
        b : in Bit;
        x : in Bit;
        y : in Bit;
        c : out Bit);
end entity clb
```

you are to devise a CLB (configurable logic block) which can be configured in order to implement the logical gates and, or, nand, and nor. More precisely, the output c of the above CLB should be:

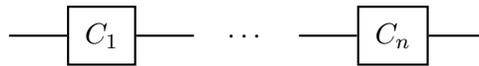
$$c = \begin{cases} a \wedge b, & \text{if } x = 0 \text{ and } y = 0, \\ a \vee b, & \text{if } x = 0 \text{ and } y = 1, \\ \neg(a \wedge b), & \text{if } x = 1 \text{ and } y = 0, \\ \neg(a \vee b), & \text{if } x = 1 \text{ and } y = 1. \end{cases}$$

- (a) Write a *behavioral* VHDL architecture for the CLB.
- (b) Draw a hardware schematic, *whose basic components are the logical gates not, and, and or.*
- (c) Assuming that each gate has a WCET of 10ns, what is the WCET of your CLB's structural architecture?

Problem 5 (Reliability)

[10 points]

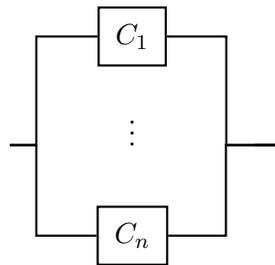
If n independent modules with reliability $R(t) = e^{-\lambda t}$ are composed sequentially, what is the mean time to failure (MTTF) of the whole system?



Problem 6 (Reliability)

[10 points]

You are to build a system S by composing in parallel n components C_1, \dots, C_n .

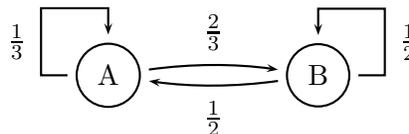


Each component C_i has reliability $\frac{1}{2}$ and costs 100 euros. What is the minimal cost of system S , if its reliability is at least $\frac{9}{10}$?

Problem 7 (Markov processes)

[10 points]

Given the following Markov process, compute $s_{\lim}(A)$ and $s_{\lim}(B)$.



Problem 8 (A/D conversion)**[10 points]**

Consider the problem of converting an analog value ranging between $U_{\min} = 3V$ and $U_{\max} = 18.5V$ into a corresponding 5-bit digital value ranging between 00000_2 and 11111_2 . Using the successive approximation method, carry out the conversion of the input voltages $U_{\text{in}} = 9.2V$ and $16.4V$ into the corresponding binary values. In each case, and for each step of the conversion, show the arranged comparison voltage U_{ref} as well as the binary value after each comparison.

Problem 9 (Scheduling)**[10 points]**

Let $\mathcal{J} = \{J_1, \dots, J_5\}$ be an aperiodic synchronous task set with

$$\begin{array}{ll} C_1 = 6 & d_1 = 3 \\ C_2 = 4 & d_2 = 8 \\ C_3 = 4 & d_3 = 13 \\ C_4 = 3 & d_4 = 16 \\ C_5 = 9 & d_5 = 24 \end{array}$$

Construct a schedule σ_{\max} that minimizes the maximum lateness L_{\max} , as well as a schedule σ_{late} that minimizes the number of late tasks N_{late} .

Problem 10 (Scheduling)**[10 points]**

Construct a periodic task set Γ with the following properties:

1. each task τ_i has phase $\Phi_i = 0$;
2. Γ is schedulable by EDF;
3. Γ is not schedulable by RM;

For each task τ_i , specify the period T_i , the relative deadline D_i , and the computation time C_i .

Problem 11 (Scheduling)**[10 points]**

Construct a periodic task set Γ with the following properties:

1. each task τ_i has phase $\Phi_i = 0$;
2. Γ is schedulable by DM (Deadline Monotonic);
3. Γ is not schedulable by RM (Rate Monotonic).

For each task τ_i , specify the period T_i , the relative deadline D_i , and the computation time C_i .

Problem 12 (Scheduling)

[10 points]

Construct a periodic task set Γ with the following properties:

1. each task τ_i has phase $\Phi_i = 0$;
2. $U(\Gamma) = \frac{1}{2}$;
3. Γ is not schedulable by EDD.

For each task τ_i , specify the period T_i , the relative deadline D_i , and the computation time C_i .

Problem 13 (Scheduling)

[20 points]

Consider the problem of scheduling a set of synchronous aperiodic tasks on a uniprocessor machine. The algorithm SJF (shortest job first) executes the tasks in nondecreasing order of computation time.

For each of the following criteria, determine whether the SJF algorithm minimizes it:

- (a) maximum lateness $L_{max} = \max_i(f_i - d_i)$;
- (b) average response time $\bar{R} = \frac{1}{n} \sum_{i=1}^n f_i$;
- (c) total completion time $T_c = \sum_{i=1}^n f_i$;
- (d) weighted sum of completion times $T_w = \sum_{i=1}^n w_i f_i$;
- (e) number of late tasks $N_{late} = \sum_{i=1}^n (if f_i > d_i then 1 else 0)$.

For each criterion, if SJF minimizes it, give a formal proof; otherwise give a counterexample.